

Dissecting Google's AddressSanitizer (ASAN)

Vishal Juneja, Computer Science (Cybersecurity)

Mentor: Yan Shoshitaishvili, Associate Professor, School of Computing and Augmented Intelligence



Introduction

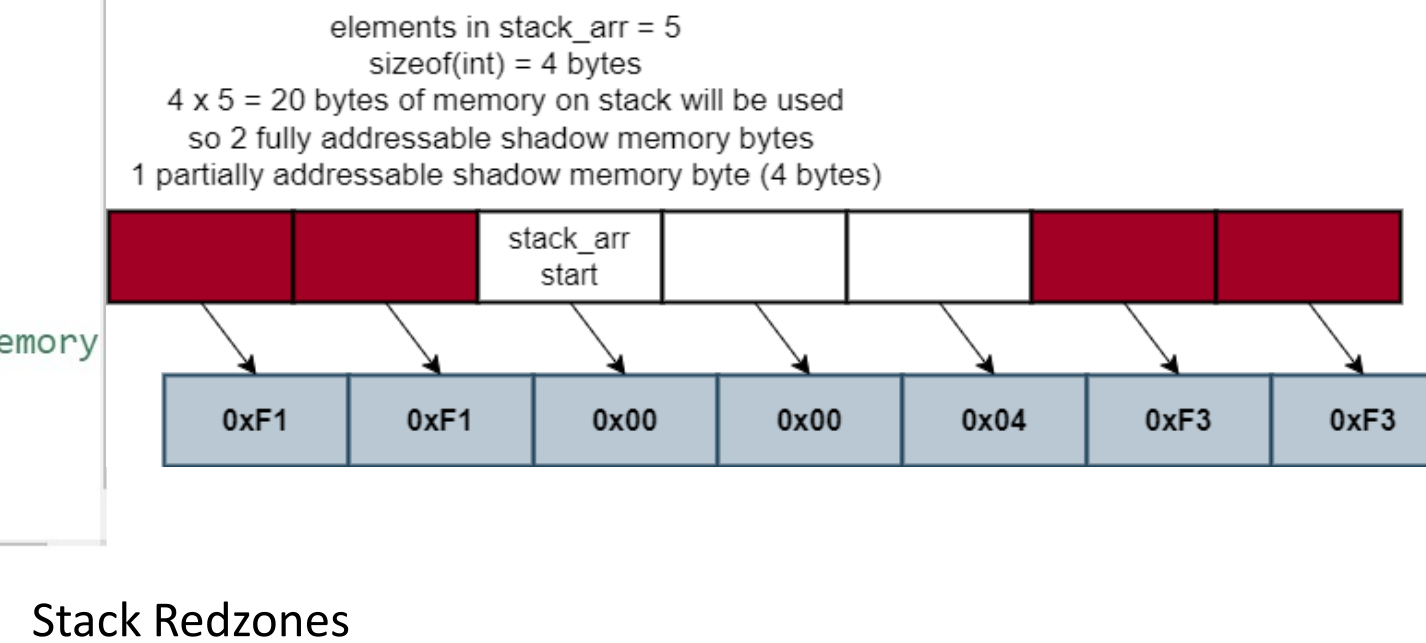
AddressSanitizer (ASAN) is an automatic error detection tool made by Google. It is used to find bugs such as Use-After-Free, Out-Of-Bound accesses to heap, stack and global objects etc.

Observations

1. We researched on the inner workings of ASAN and identified the important components.
2. The backbone of ASAN involves shadow memory, which stores metadata about the program's memory usage.
3. Shadow memory indicates whether memory is accessible and its state, such as if it is accessible, partially accessible, or inaccessible.
4. The correspondence between program memory and shadow memory is one shadow memory byte for every 8 bytes of program memory.
5. A shadow memory byte of 0 indicates that all 8 bytes are accessible, while other values from 1 to 7 represent partial accessibility, and -1 means no memory is accessible.
6. The ASAN Shadow Memory address calculation formula allows the computation of shadow memory addresses.
7. Instrumentation involves checking the shadow memory byte before memory access, introducing a time overhead for enhanced security.
8. The runtime library in ASAN manages shadow memory, replacing standard memory allocation functions with specialized implementations.
9. Malloc allocates redzones around memory regions to detect overflows and underflows, with redzone size influencing detection scope.
10. Free function poisons memory regions and places them in quarantine as a FIFO queue.
11. Redzones for global objects are generated at compile time.

```
#include <stdio.h>
#include <string.h>

int main(int argc, char *argv[])
{
    int stack_arr[5];
    stack_arr[6] = 2; // accessing out-of-bound memory
    return 0;
}
```



```
SUMMARY: AddressSanitizer: heap-buffer-overflow (/home/kinono/research/test/sample_malloc_with_asan+0x4998a6) (BuildId: e680c6dd6d724ab83a7a4d63a446693a65fed3a) in strcpy
Shadow bytes around the buggy address:
0x0c077fff7fde: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0x0c077fff7fd0: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0x0c077fff7fe8: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0x0c077fff8000: fa fa 00[02]fa fa fa fa fa fa fa fa fa fa fa fa
0x0c077fff8010: fa fa fa fa fa fa fa fa fa fa fa fa fa fa fa fa
0x0c077fff8020: fa fa fa fa fa fa fa fa fa fa fa fa fa fa fa fa
0x0c077fff8030: fa fa fa fa fa fa fa fa fa fa fa fa fa fa fa fa
0x0c077fff8040: fa fa fa fa fa fa fa fa fa fa fa fa fa fa fa fa
0x0c077fff8050: fa fa fa fa fa fa fa fa fa fa fa fa fa fa fa fa
Shadow byte Legend (one shadow byte represents 8 application bytes):
Addressable: 00
Partially addressable: 01 02 03 04 05 06 07
Heap left redzone: fa
Fixed heap redzone: fd
Stack left redzone: f1
Stack mid redzone: f2
Stack right redzone: f3
Stack after return: f5
Stack use after scope: f8
Global init order: f9
Poisoned by user: f7
Container overflow: fc
Array cookie: ac
Intra object redzone: bb
ASAN internal: fe
Left alloca redzone: ca
Right alloca redzone: cb
==2799==ABORTING
```

Heap overflow when ASAN is enabled

```
.text:00000000004D9EA0 ; char __fastcall load(char *addr)
.text:00000000004D9EA0 public load
.text:00000000004D9EA0 load ; CODE XREF: main+CB4p
.text:00000000004D9EA0 var_18 = qword ptr -18h
.text:00000000004D9EA0 var_9 = byte ptr -9
.text:00000000004D9EA0 var_8 = qword ptr -8
.text:00000000004D9EA0 addr = qword ptr 8
.text:00000000004D9EA0 ; __unwind {
.text:00000000004D9EA0 push rbp
.text:00000000004D9EA1 mov rbp, rsp
.text:00000000004D9EA4 sub rsp, 20h
.text:00000000004D9EA8 mov [rbp+var_8], rdi
.text:00000000004D9EAC mov rax, [rbp+var_8]
.text:00000000004D9EB0 mov [rbp+var_18], rax
.text:00000000004D9EB4 shr rax, 3
.text:00000000004D9EB8 mov al, [rax+7FFF8000h]
.text:00000000004D9EBE mov [rbp+var_9], al
.text:00000000004D9EC1 cmp al, 0
.text:00000000004D9EC3 jz loc_4D9EE5
.text:00000000004D9EC9 mov cl, [rbp+var_9]
.text:00000000004D9ECC mov rax, [rbp+var_18]
.text:00000000004D9ED0 and rax, 7
.text:00000000004D9ED4 cmp al, cl
.text:00000000004D9ED6 jl loc_4D9EE5
.text:00000000004D9EDC mov rdi, [rbp+var_18] ; unsigned __int64
.text:00000000004D9EE0 call __asan_report_load1
.text:00000000004D9EE5 loc_4D9EE5: ; CODE XREF: load+23tj
.text:00000000004D9EE5 ; load+36tj
.text:00000000004D9EE5 mov rax, [rbp+var_18]
.text:00000000004D9EE9 movsx eax, byte ptr [rax]
.text:00000000004D9EF0 add rsp, 20h
.text:00000000004D9EF1 pop rbp
.text:00000000004D9EF1 retn
```

Checks Added by instrumentation module

Problems

1. ASAN is not very well documented. Current documentation is very concise, not very user friendly and lacks detailed explanation. On top of it, resources for learning ASAN is very scattered. The error messages produced by ASAN are difficult to interpret.
2. ASAN has high memory and time overhead. Average slowdown caused by ASAN is 73% [1].

Future Work

Currently, we are conducting tests on various programs, calculating the overall memory and time overhead caused by ASAN. We are in the process of identifying the components that are causing the overhead, and then we plan to enhance the performance of these components by improving their functionality.

Background Study

1. K. Serebryany, D. Bruening, A. Potapenko, and D. Vyukov, "AddressSanitizer: A Fast Address Sanity Checker," in 2012 USENIX Annual Technical Conference (USENIX ATC 12), Jun. 2012, pp. 309–318. [Online]. Available: <https://www.usenix.org/conference/atc12/technical-sessions/presentation/serebryany>
2. Y. Zhang, C. Pang, G. Portokalidis, N. Triandopoulos, and J. Xu, "Debloating Address Sanitizer," in 31st USENIX Security Symposium (USENIX Security 22), Aug. 2022, pp. 4345–4363. [Online]. Available: <https://www.usenix.org/conference/usenixsecurity22/presentation/zhang-yuchen>
3. <https://github.com/gcc-mirror/gcc/blob/releases/gcc-12.2.0/gcc/asan.cc>